Data on the Heap

num
bool
x char
x double

Next, lets add support for

- **Data Structures**

In the process of doing so, we will learn about

- **Heap Allocation**
- **Run-time Tags**
- High-order Func (Closures)

<env, code>

Creating Heap Data Structures

We have already support for *two* primitive data types

```
data Ty
  = TNumber       -- e.g. 0,1,2,3,...
  | TBoolean      -- e.g. true, false
```

we could add several more of course, e.g.

- `Char`
- `Double` or `Float`

etc. (you should do it!)

However, for all of those, the same principle applies, more or less

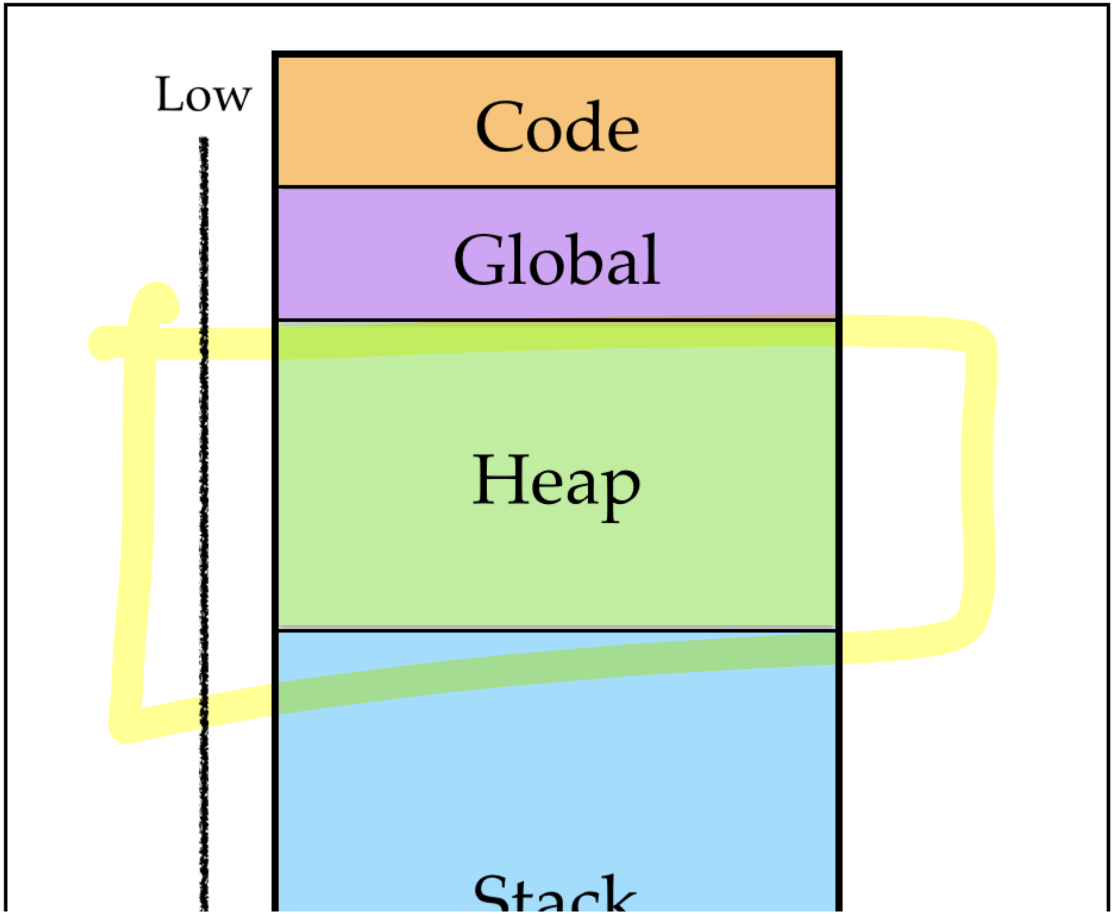
- As long as the data fits into a single word (8-bytes)

Instead, lets learn how to make **unbounded data structures**

- Lists
- Trees
- ...

which require us to put data on the **heap**

> *not just the stack that we've used so far.*

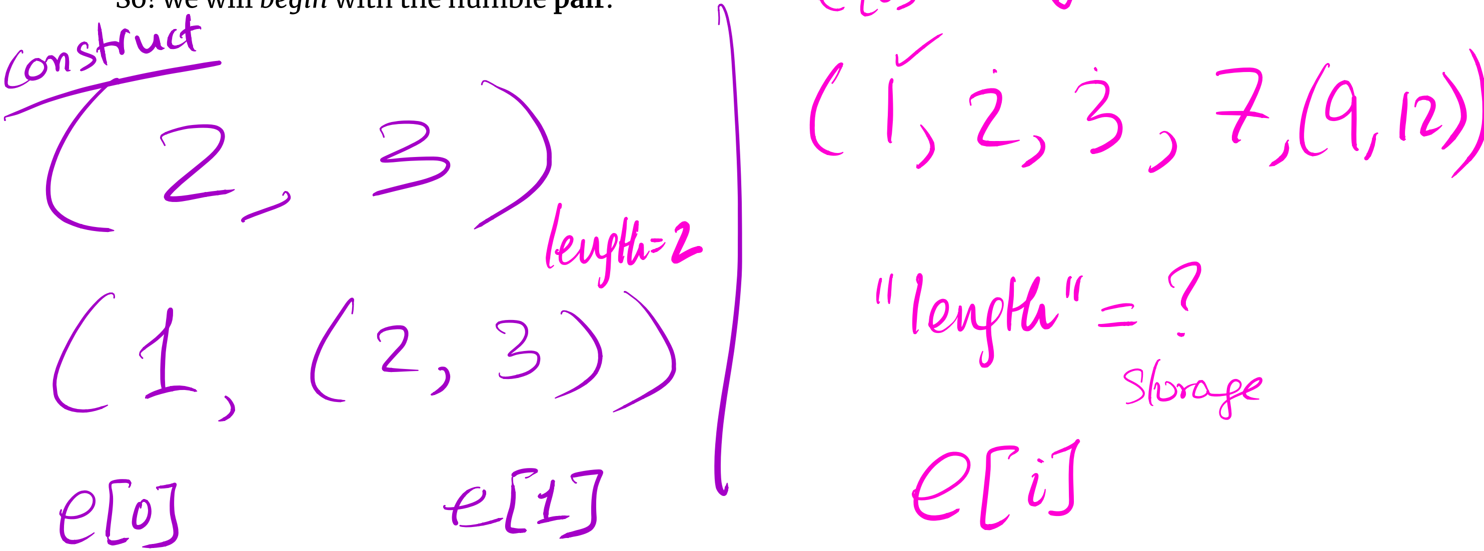


Stack vs. Heap

Pairs

While our *goal* is to get to lists and trees, the journey of a thousand miles begins with a single step…

So! we will *begin* with the humble **pair**.



Pairs: Semantics (Behavior)

First, lets ponder what exactly we're trying to achieve.

We want to enrich our language with *two* new constructs:

- **Constructing** pairs, with a new expression of the form `(e0, e1)` where `e0` and `e1` are expressions.
- **Accessing** pairs, with new expressions of the form `e[0]` and `e[1]` which

evaluate to the first and second element of the tuple `e` respectively.

For example,

```
let t = (2, 3) in
  t[0] + t[1]
```

should evaluate to `5`.

Strategy

Next, lets informally develop a strategy for extending our language with pairs, implementing the above semantics. We need to work out strategies for:

1. **Representing** pairs in the machine's memory,

$(e_0, e_1) \longrightarrow$ <asm>

2. **Constructing** pairs (i.e. implementing `(e0, e1)` in assembly),
3. **Accessing** pairs (i.e. implementing `e[0]` and `e[1]` in assembly).

1. Representation

Recall that we represent all values: (05-cobra.md/#option-2-use-a-tag-bit)

- `Number` like `0`, `1`, `2` ...
- `Boolean` like `true`, `false`

as a **single word** either

- 8 bytes on the stack, or
- a single register `rax`, `rbx` etc.

EXERCISE

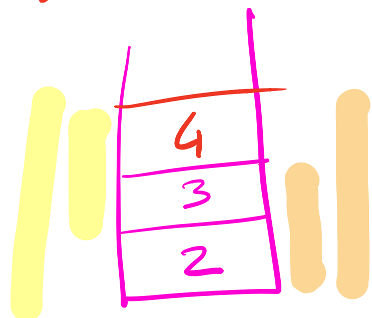
What kinds of problems do you think might arise if we represent a pair `(2, 3)` on the *stack* as:

```
|     |
-------
|  3  |
-------
|  2  |
-------
| ... |
-------
```

let t = ((2, 3), 4) in

t = ((2, 3), 4)

t = (2, (3, 4))



t[0][0]

1, 2, 3, 4, 5

cons(1, cons(2, cons(3, cons(4, nil()))))

(1, (2, (3, (4, false))))

(1,2) 3 → (e0, e1)

e[0] , e[1]

```
def tail(l):
    l[1]

def isNil(l):
    l == False

def nil():
    false

def cons(h,t):
    (h,t)

def range(lo, hi):
    if lo<hi:
        cons(lo, range(lo+1, hi))
    else:
        nil()

def length(l):
    if isNil(l):
        0
    else:
        1+length(tail(l))

l = range(0,100)
length(l)
```

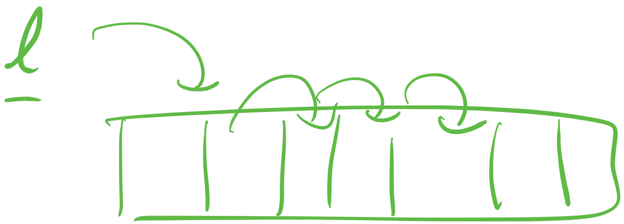
build list

QUIZ

How many words would we need to store the tuple

(3, (4, 5))

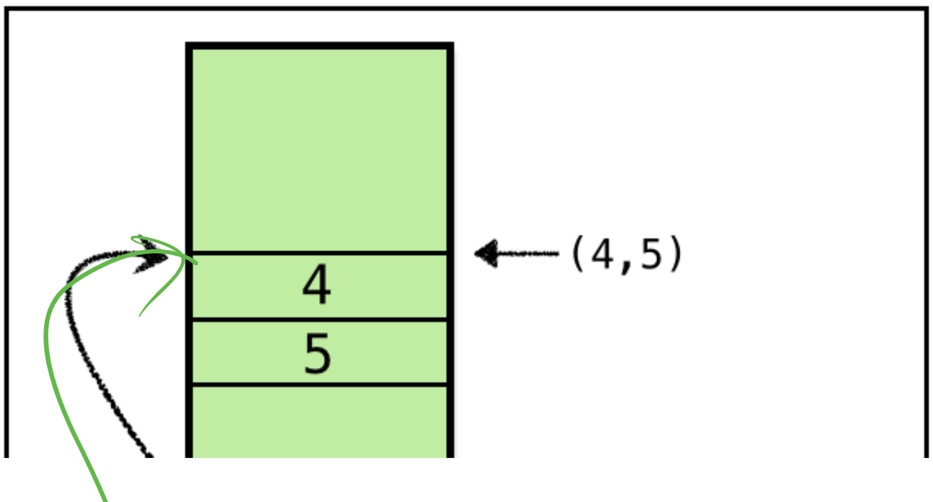
1. ~~1 word~~
2. 2 words
3. 3 words
4. 4 words
5. 5 words

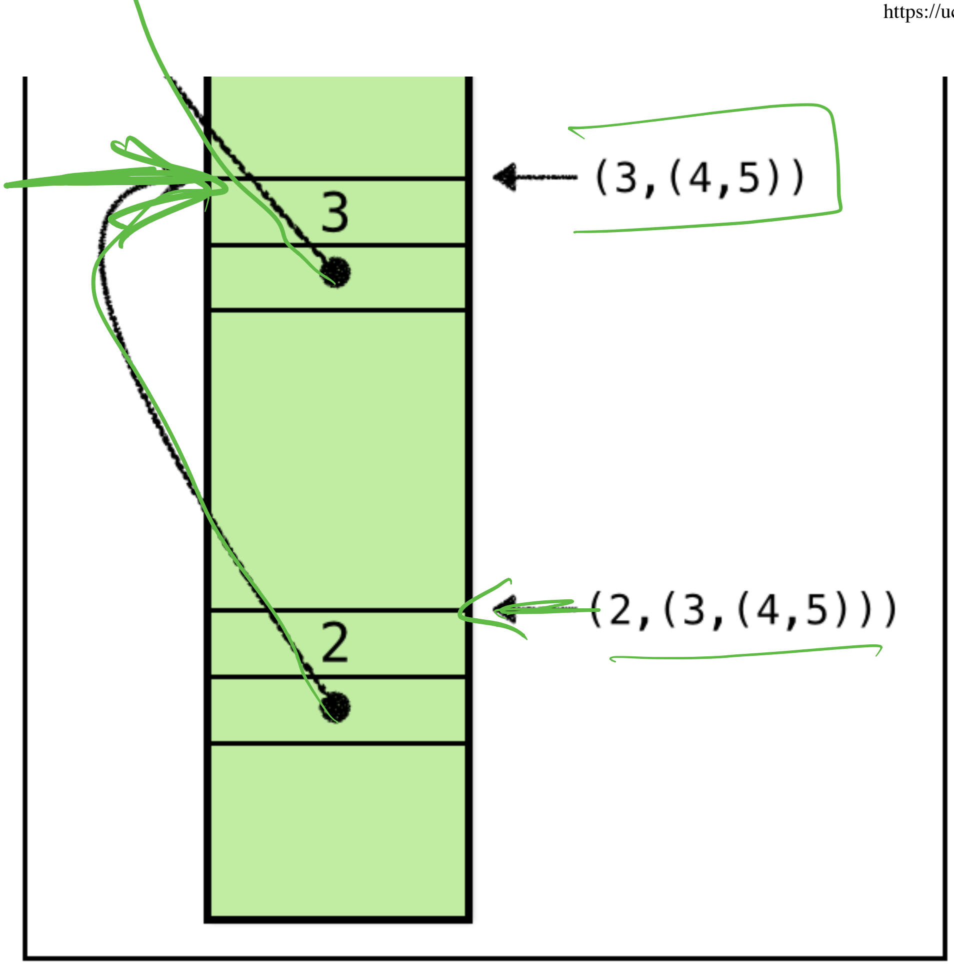


Pointers

> *Every problem in computing can be solved by adding a level of indirection.*

We will **represent a pair** by a **pointer** to a block of **two adjacent words** of memory.





Pairs on the heap

The above shows how the pair `(2, (3, (4, 5)))` and its sub-pairs can be stored in the **heap** using pointers.

`(4, 5)` is stored by adjacent words storing

- `4` and
- `5`

`(3, (4, 5))` is stored by adjacent words storing

- `3` and
- a **pointer** to a heap location storing `(4, 5)`

`(2, (3, (4, 5)))` is stored by adjacent words storing

- `2` and
- a **pointer** to a heap location storing `(3, (4, 5))`.

A Problem: Numbers vs. Pointers?

How will we tell the difference between *numbers* and *pointers*?

That is, how can we tell the difference between

1. the *number* 5 and
2. a *pointer* to a block of memory (with address `5`)?

Each of the above corresponds to a *different* tuple

1. `(4, 5)` or
2. `(4, (...))`.

so its pretty crucial that we have a way of knowing *which* value it is.

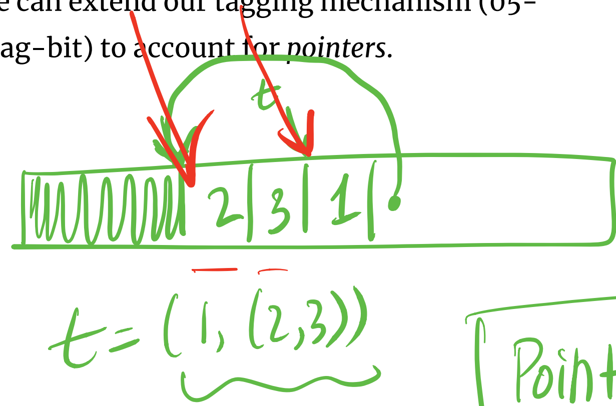
t = (1, (2, 3))

t = 3

t = false

Tagging Pointers

As you might have guessed, we can extend our tagging mechanism (05-cobra.md/#option-2-use-a-tag-bit) to account for *pointers*.



Type	LSB
`number`	xx0
`boolean`	111
`pointer`	001

That is, for

- `number` the **last bit** will be `0` (as before),
- `boolean` the **last 3 bits** will be `111` (as before), and
- `pointer` the **last 3 bits** will be `001`.

(We have 3-bits worth for tags, so have wiggle room for other primitive types.)

Address Alignment

As we have a **3 bit tag**

- leaving **64 - 3 = 61 bits** for the actual address

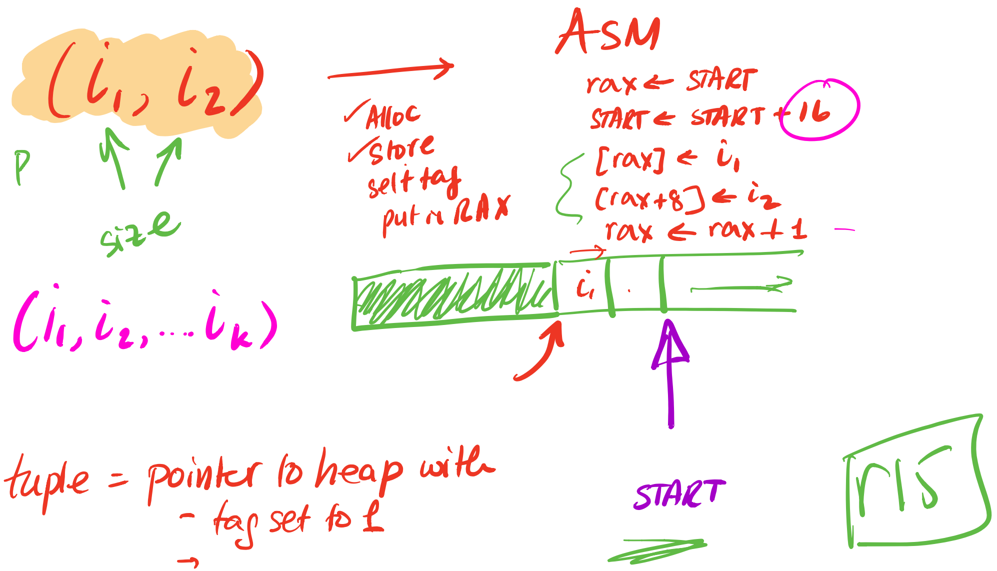
So actual addresses, written in binary, omitting trailing zeros, are of the form

Binary	Decimal
`0b00000000`	0
`0b00001000`	8
`0b00010000`	16
`0b00011000`	24
`0b00100000`	32
...	

That is, the addresses are **8-byte aligned**.

Which is great because at each address, we have a pair, i.e. a **2-word = 16-byte block**, so the *next* allocated address will *also* fall on an 8-byte boundary.

- But … what if we had 3-tuples? or 5-tuples? …



2. Construction

Next, lets look at how to implement pair **construction** that is, generate the assembly for expressions like:

```
(e1, e2)
```

To **construct** a pair `(e1, e2)` we

1. **Allocate** a new 2-word block, and getting the starting address at `rax`,
2. **Copy** the value of `e1` (resp. `e2`) into `[rax]` (resp. `[rax + 8]`).

3. **Tag** the last bit of `rax` with `1`.

The resulting `eax` is the **value of the pair**

- The *last step* ensures that the value carries the proper tag.

ANF will ensure that `e1` and `e2` are immediate expressions (04-boa.md/#idea-immediate-expressions)

- will make the second step above straightforward.

EXERCISE How will we do ANF conversion for `(e1, e2)`?

Allocating Addresses

Lets use a **global** register `r15` to maintain the address of the **next free block** on the heap.

Every time we need a *new* block, we will:

1. Copy the current `r15` into `rax`

- Set the last bit to `1` to ensure proper tagging.
- `rax` will be used to fill in the values

2. Increment the value of `r15` by `16`

- Thus *allocating* 8 bytes (= 2 words) at the address in `rax`

Note that addresses stay 8-byte aligned (last 3 bits = 0) if we

- *Start* our blocks at an 8-byte boundary, and
- *Allocate* 16 bytes at a time,

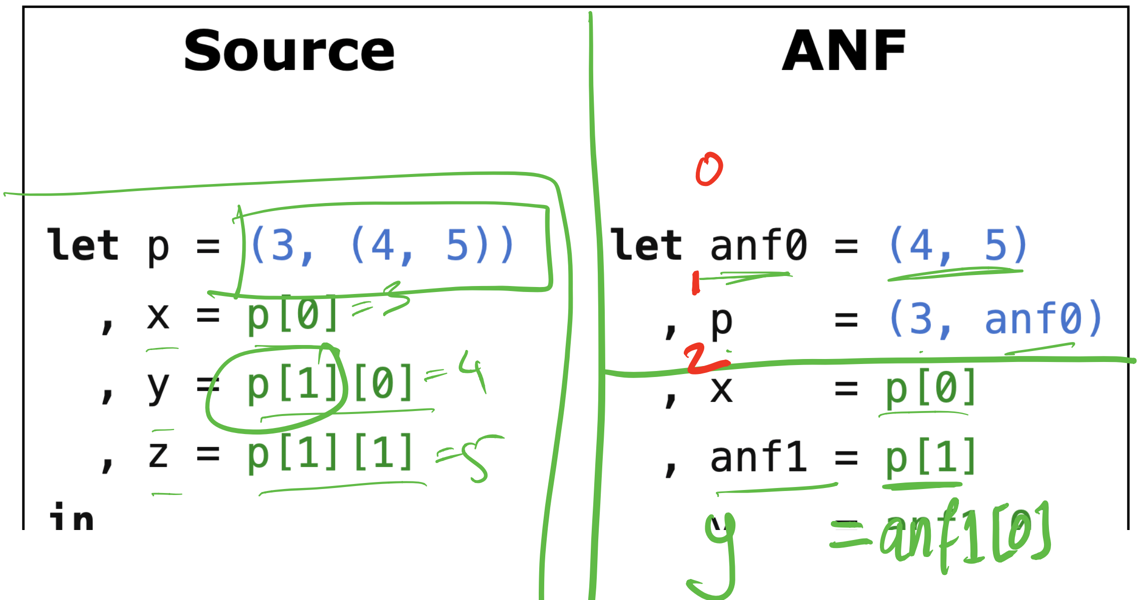
NOTE: Your assignment will have *blocks of varying sizes*

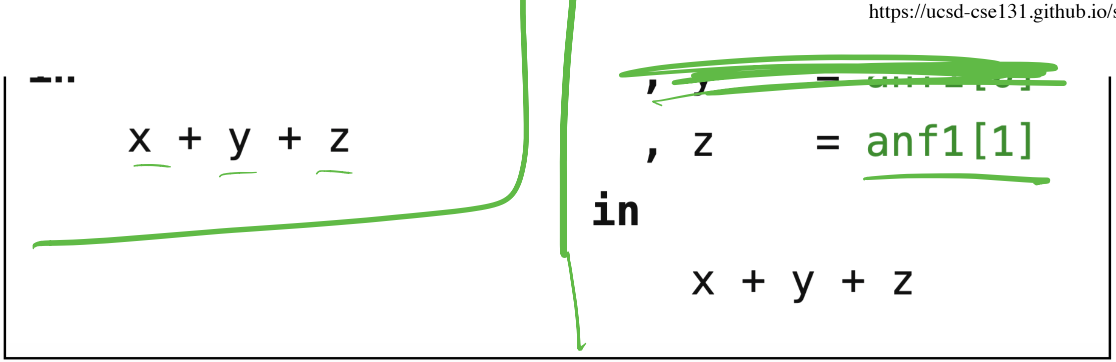
- You will have to *maintain* the 8-byte alignment by *padding*

Example: Allocation

In the figure below, we have

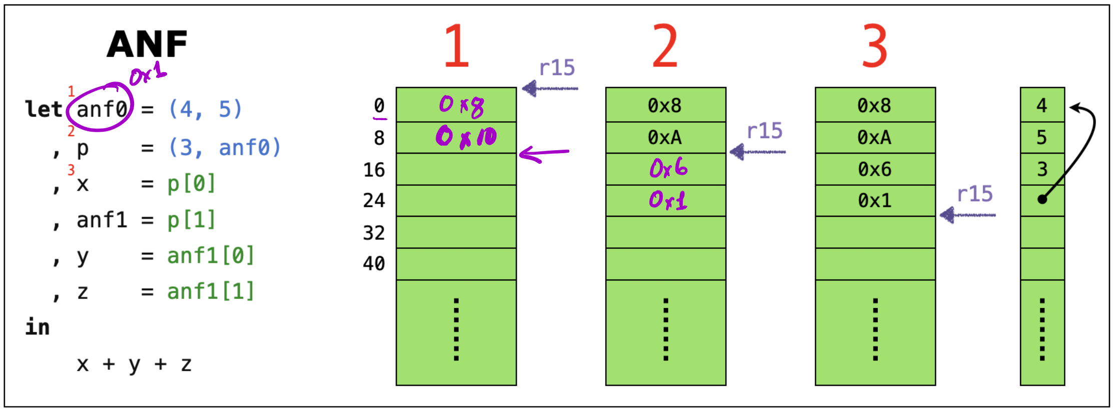
- a source program on the left,
- the ANF equivalent next to it.





Example of Pairs

The figure below shows the how the heap and `r15` evolve at points 1, 2 and 3:

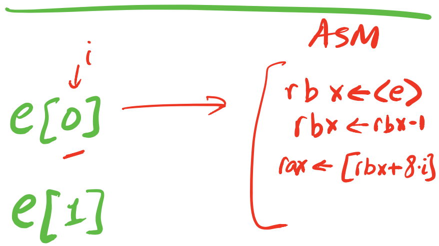


Allocating Pairs on the Heap

QUIZ

In the ANF version, `p` is the *second (local) variable* stored in the stack frame. What *value* gets moved into the *second stack slot* when evaluating the above program?

1. `0x3` → 3
2. `(3, (4, 5))` →
3. `0x11` → 17
4. `0x9` → 9
5. `0x10` → 16



3. *Accessing*

Finally, to **access** the elements of a pair

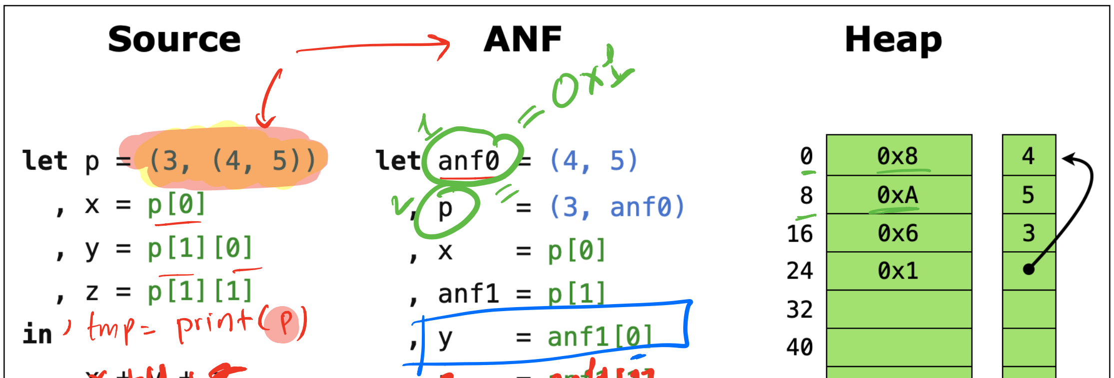
Lets compile `e[0]` to get the first or `e[1]` to get the second element

1. **Check** that immediate value `e` is a pointer
2. **Load** `e` into `rbx`
3. **Remove** the tag bit from `rbx`

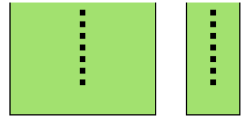
4. **Copy** the value in `[rbx]` (resp. `[rbx + 8]`) into `rbx`.

Example: Access

Here is a snapshot of the heap after the pair(s) are allocated.



```
    in
      x + y + z
```



Allocating Pairs on the Heap

Lets work out how the values corresponding to `x`, `y` and `z` in the example above get stored on the stack frame in the course of evaluation.

Variable	Hex Value	Value
`anf0`	`0x001`	`ptr 0`
`p`	`0x011`	`ptr 16`
`x`	`0x006`	`num 3`
`anf1`	`0x001`	`ptr 0`
`y`	`0x008`	`num 4`
`z`	`0x00A`	`num 5`
`anf2`	`0x00E`	`num 7`
result	`0x018`	`num 12`

Plan

Pretty pictures are well and good, time to build stuff!

As usual, lets continue with our recipe:

1. Run-time
2. Types
3. Transforms

We've already built up intuition of the *strategy* for implementing tuples. Next, lets look at how to implement each of the above.

Run-Time

We need to extend the run-time (`c-bits/main.c`) in two ways.

1. **Allocate** a chunk of space on the heap and pass in start address to `our_code` .
2. **Print** pairs properly.

Allocation

The first step is quite easy we can use `calloc` as follows:

```
int main(int argc, char** argv) {
  int* HEAP = calloc(HEAP_SIZE, sizeof (int));
  long result = our_code_starts_here(HEAP);
  print(result);
  return 0;
}
```

where does 'HEAP' live? in our_code....

The above code, (A) r15 (B) rdi

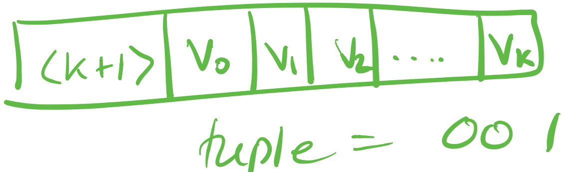
1. **Allocates** a big block of contiguous memory (starting at `HEAP`), and
2. **Passes** this address in to `our_code`.

Now, `our_code` needs to, at the beginning start with instructions that

- copy the parameter (in `rdi`) into global pointer (`r15`) (8-byte?)
- and then bump it up at each allocation.

→ alloc

→ print ✓



Printing

To print pairs, we must **recursively traverse** pointers

- until we hit `number` or `boolean`.

We can check if a value is a pair by looking at its last 3 bits:

```
int isPair(int p) {
  return (p & 0x00000007) == 0x00000001;
}
```

We can use the above test to recursively print (word)-values:

```
void print(long val) {
  if(val & 0x1 == 0) { // val is a number
    printf("%ld", val >> 1);
  }
  else if(val == CONST_TRUE) {        // val is true
    printf("true");
  }
  else if(val == CONST_FALSE) {       // val is false
    printf("false");
  }
  else if(val & 7 == 1) {
    long* valp = (long *) (val - 1);  // extract address
    printf("(");
    print(*valp);                     // print first element
    printf(", ");
    print(*(valp + 1));               // print second element
    printf(")");
  }
  else {
    printf("Unknown value: %#010x", val);
  }
}
```

$e_1 [e_2]$

Types

Next, lets move into our compiler, and see how the **core types** need to be extended.

Source

$(e_1, e_2) \longrightarrow$ Pair e_1 e_2

We need to extend the *source* `Expr` with support for tuples

e[0] $\longrightarrow$ GetItem e First

```
data Expr a
  = ...
  | Pair    (Expr a) (Expr a) a   -- ^ construct a pair
  | GetItem (Expr a) Field    a   -- ^ access a pair's element
```

In the above, `Field` is

e[1] $\longrightarrow$ GetItem e Second

| Tuple [Expr a] a

| GetItem (Expr a) Int → "static"

(Expr a) → "dynamic"

```
data Field
  = First      -- ^ access first element of pair
  | Second     -- ^ access second element of pair
```

NOTE: Your assignment will generalize pairs to **n-ary tuples** using

- `Tuple [Expr a]` representing `(e1,...,en)`
- `GetItem (Expr a) (Expr a)` representing `e1[e2]`

Dynamic Types

Let us extend our **dynamic types** `Ty` see (05-cobra.md/#types) to include pairs:

```
data Ty = TNumber | TBoolean | TPair
```

0x0 0x111 0x001

Assembly

The assembly `Instruction` are changed minimally; we just need access to `r15` which will hold the value of the *next* available memory block:

```
data Register
  = ...
  | R15
```

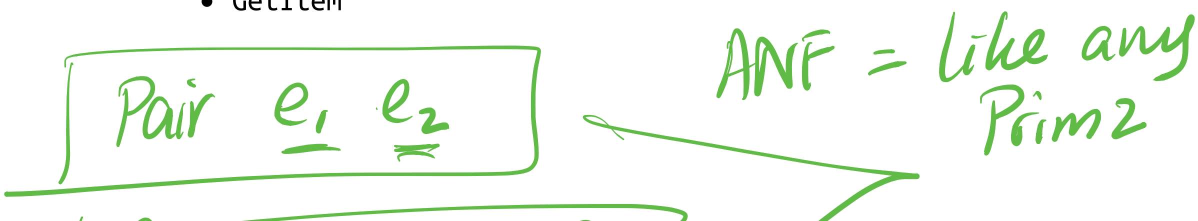
Transforms

Our code must take care of three things:

1. **Initialize** `r15` to allow heap allocation,
2. **Construct** pairs,
3. **Access** pairs.

The latter two will be pointed out as cases in `anf` and `compileEnv`

- `Tuple`
- `GetItem`



Initialize

We need to **initialize** `r15` with the **start position** of the heap

- passed in as `rdi` by the run-time.

How shall we get a hold of this position?

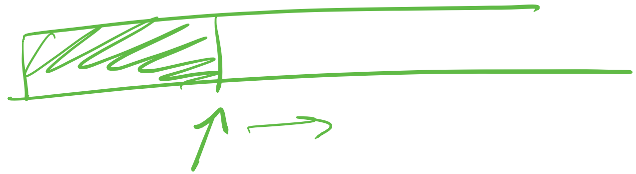
To do so, `our_code` starts off with a `prelude`

```
prelude :: [Instruction]
prelude =
  [ IMov (Reg R15) (Reg RDI)   -- copy param (HEAP) off rdi
  ]
```

Is that it?

(1) find the gap
add the gap as pad

(2) zero out last bits



QUIZ

Is `r15` 8-byte aligned?

A. Yes

B. No ✗

Ensuring alignment


```
prelude :: [Instruction]
prelude =
  [ IMov (Reg RAX) (HexConst 0xFFFFFFFF)      -- setup regMask
  , IShl (Reg RAX) (Const 32)
  , IOr  (Reg RAX) (HexConst 0xFFFFFFF8)
  , IMov (Reg R15) (Reg RDI)                  -- copy param (HEAP) of
f rdi
  , IAdd (Reg R15) (Const 8)                  -- add 8 and mask 3 bit
s to ensure
  , IAnd (Reg R15) (Reg RAX)                  -- 8-byte aligned
  ]
```

1. **Copy** the value off the (parameter) stack, and
2. **Adjust** the value to ensure the value is *8-byte aligned.*

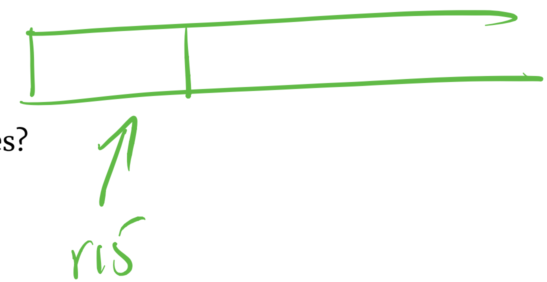
```
mov r15, rdi
add r15, 8
and r15, 0x111....000
```

61-bits

QUIZ

Why add `8 to r15`? What would happen if we *removed* that operation?

1. `r15` would not be 8-byte aligned?
2. `r15` would point into the stack?
3. `r15` would not point into the heap?
4. `r15` would not have enough space to write 2 bytes?



Construct

To *construct* a pair `(v1, v2)` we directly implement the above strategy (07-egg-eater.md/#2-construction):

```
compileEnv env (Tuple v1 v2)
  = pairAlloc                              -- 1. allocate pair, resulting addr in `rax`
 ++ pairCopy First  (immArg env v1)        -- 2. copy first value into slot
 ++ pairCopy Second (immArg env v2)        -- 3. copy second value into slot
 ++ setTag   RAX      TPair                -- 3. set the tag-bits of `rax`
```

Lets look at each step in turn.

```
(Tuple vs)        where n = length vs
tupleAlloc (n+1)
add RAX, 1

pairAlloc =
    mov rax, r15
    add r15, 16

tupleAlloc k =
    mov rax, r15
    add r15, 8·k

pairCopy fld arg =
    mov rbx, arg
    mov [rax+off], rbx
  where
    off = fieldOff fld

[rax + off]
[rbp ± _]
```

Allocate

r15

To allocate, we just copy the current pointer `r15` and increment by `16` bytes,

- accounting for two 8-byte blocks for each element.

```
pairAlloc :: Asm
pairAlloc
  = [ IMov (Reg RAX) (Reg R15)    -- copy current "free address" `esi
` into `eax`
    , IAdd (Reg RAX) (Const 16)   -- increment `esi` by 8
    ]
```

Exercise How would you make this work for `n` -tuples?

Copy

We copy an `Arg` into a `Field` by

- saving the `Arg` into a helper register `rbx`,
- copying `rbx` into the field's slot on the heap.

```
pairCopy :: Field -> Arg -> Asm
pairCopy fld arg
  = [ IMov (Reg RBX) arg
    , IMov (pairAddr fld) (Reg RBX)
    ]
```

Recall, the field's slot is either `[rax]` or `[rax + 8]` depending on whether the field is `First` or `Second`.

QUIZ

What shall we fill in for `_1` and `_2` ?

```
pairAddr :: Field -> Arg
pairAddr First  = RegOffset ?1 RAX
pairAddr Second = RegOffset ?2 RAX
```

A. `0 and 1`

B. `0 and -1`

C. `1` and `2`

D. `-1` and `-2`

E. huh?

Tag

Finally, we set the tag bits of `rax` by using `typeTag TPair` which is defined

```
setTag :: Register -> Asm
setTag r = [ IAdd (Reg r) (HexConst 0x1) ]
```

e[0] e[1]

Access

To access tuples, lets update `compileEnv` with the strategy above:

```
compileExpr env (GetItem e fld)
  = assertType env e TPair                   -- 1. check that e is a (pai
r) pointer
 ++ [ IMov (Reg RAX) (immArg env e) ]        -- 2. load pointer into eax
 ++ unsetTag RAX                             -- 3. remove tag bit to get a
ddress
 ++ [ IMov (Reg RAX) (pairAddr fld) ]        -- 4. copy value from resp. s
lot to eax
```

we remove the tag bits by doing the opposite of `setTag` namely:

```
unsetTag :: Register -> Asm
unsetTag r = ISub (Reg RAX) (HexConst 0x1)
```

N-ary Tuples

Thats it! Lets take our compiler out for a spin, by using it to write some interesting programs!

First, lets see how to generalize pairs to allow for

- triples `(e1,e2,e3)`
- quadruples `(e1,e2,e3,e4)`
- pentuples `(e1,e2,e3,e4,e5)`

(e1,(e2,e3))

(e1,(e2,(e3,(e4,_))))

and so on.

We just need a library of functions in our new `egg` language to

- **Construct** such tuples, and
- **Access** their fields.

Constructing Tuples

We can write a small set of functions to **construct** tuples (up to some given size):

```
def tup3(x1, x2, x3):
  (x1, (x2, x3))

def tup4(x1, x2, x3, x4):
  (x1, (x2, (x3, x4)))

def tup5(x1, x2, x3, x4, x5):
  (x1, (x2, (x3, (x4, x5))))
```

Accessing Tuples

We can write a single function to access tuples of any size.

So the below code

```
let yuple = (10, (20, (30, (40, (50, false))))) in

get(yuple, 0) = 10
get(yuple, 1) = 20
get(yuple, 2) = 30
get(yuple, 3) = 40
get(yuple, 4) = 50
```

```
def tup3(x1, x2, x3):
  (x1, (x2, x3))

def tup5(x1, x2, x3, x4, x5):
  (x1, (x2, (x3, (x4, x5))))

let t  = tup5(1, 2, 3, 4, 5) in
  , x0 = print(get(t, 0))
  , x1 = print(get(t, 1))
  , x2 = print(get(t, 2))
  , x3 = print(get(t, 3))
  , x4 = print(get(t, 4))
in
```

```
  99
```

should print out:

```
0
1
2
3
4
99
```

How shall we write it?

```
def get(t, i):
  TODO-IN-CLASS
```

QUIZ

Using the above "library" we can write code like:

```
let quad = tup4(1, 2, 3, 4) in
  get(quad, 0) + get(quad, 1) + get(quad, 2) + get(quad, 3)
```

What will be the result of compiling the above?

1. Compile error
2. Segmentation fault
3. Other run-time error
4. `4`
5. `10`

QUIZ

Using the above “library” we can write code like:

```
def get(t, i):
  if i == 0:
    t[0]
  else:
    get(t[1],i-1)

def tup3(x1, x2, x3):
  (x1, (x2, (x3, false)))

let quad = tup3(1, 2, 3) in
  get(quad, 0) + get(quad, 1) + get(quad, 2) + get(quad, 3)
```

What will be the result of compiling the above?

1. Compile error
2. Segmentation fault
3. Other run-time error
4. `4`
5. `10`

Lists

Once we have pairs, we can start encoding **unbounded lists**.

To build a list, we need two constructor functions:

~~~
def empty():
  false

def cons(h, t):
  (h, t)
``

We can now encode lists as:

```python
cons(1, cons(2, cons(3, cons(4, empty()))))
~~~

Access

To **access** a list, we need to know

1. Whether the list `isEmpty` , and
2. A way to access the `head` and the `tail` of a non-empty list.

```
def isEmpty(l):
  l == empty()

def head(l):
  l[0]

def tail(l):
  l[1]
```

Examples

We can now write various functions that build and operate on lists, for example, a function to generate the list of numbers between `i` and `j`

```
def range(i, j):
  if (i < j):
    cons(i, range(i+1, j))
  else:
    empty()

range(1, 5)
```

which should produce the result

```
(1,(2,(3,(4,false))))
```

and a function to sum up the elements of a list:

```
def sum(xs):
  if (isEmpty(xs)):
    0
  else:
    head(xs) + sum(tail(xs))

sum(range(1, 5))
```

which should produce the result `10`.

Recap

We have a pretty serious language now, with:

- **Data Structures**

which are implemented using

- **Heap Allocation**
- **Run-time Tags**

which required a bunch of small but subtle changes in the

- runtime and compiler

In your assignment, you will add *native* support for n-ary tuples, letting the programmer write code like:

```
(e1, e2, e3, ..., en)  # constructing tuples of arbitrary arity

e1[e2]                 # allowing expressions to be used as fields
```

Next, we'll see how to

- use the "tuple" mechanism to implement **higher-order functions** and
- reclaim unused memory via **garbage collection**.

(https://ucsd-cse131.github.io/sp21/feed.xml)

(https://twitter.com/ranjitjhala)

(https://plus.google.com/u/0/106612421534244742464)

(https://github.com/ucsd-cse131/sp21)