

# Branches and Binary Operators

## BOA: Branches and Binary Operators

Next, lets add

- Branches ( `if` -expressions) ✓
- Binary Operators ( `+`, `-`, etc.)

In the process of doing so, we will learn about

- **Intermediate Forms**
- **Normalization**

# Branches

TRUE = NON-ZERO  
FALSE = 0

Lets start first with branches (conditionals).

We will stick to our recipe of:

1. Build intuition with **examples**,
2. Model problem with **types**,
3. Implement with **type-transforming-functions**,
4. Validate with **tests**.

## *Examples*

First, lets look at some examples of what we mean by branches.

- For now, lets treat  $0$  as “false” and non-zero as “true”

## Example: If1

```

if 10:
    22
else:
    sub1(0)
  
```

→ 22

- Since 10 is not 0 we evaluate the “then” case to get 22

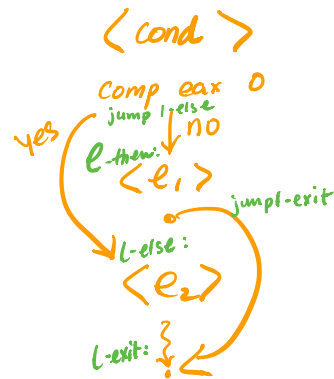
## Example: If2

```

if sub(1):
    22
else:
    sub1(0)
  
```

→ -1

if cond :  
 $e_1$   
 else :  
 $e_2$



- Since  $\text{sub}(1)$  is  $\theta$  we evaluate the “else” case to get  $-1$

## QUIZ: *If3*

`if-else` is also an *expression* so we can nest them:

What should the following evaluate to?

```
let x = if sub(1):  
        22  
        else:  
        sub1(0)  
in  
  if x:  
    add1(x)  
  else:  
    999
```

- A. 999
- B. 0
- C. 1
- D. 1000
- E. -1

# Control Flow in Assembly

To compile branches, we will use **labels**, **comparisons** and **jumps**

## Labels

our\_code\_label:

...

Labels are “landmarks” - from which execution (control-flow) can be *started*, or - to which it can be *diverted*

## *Comparisons*

```
cmp a1, a2
```

- Perform a (numeric) **comparison** between the values `a1` and `a2`, and
- Store the result in a special **processor flag**

## *Jumps*



```
jmp LABEL    # jump unconditionally (i.e. always)
je LABEL     # jump if previous comparison result was EQUAL
jne LABEL    # jump if previous comparison result was NOT-EQUAL
```

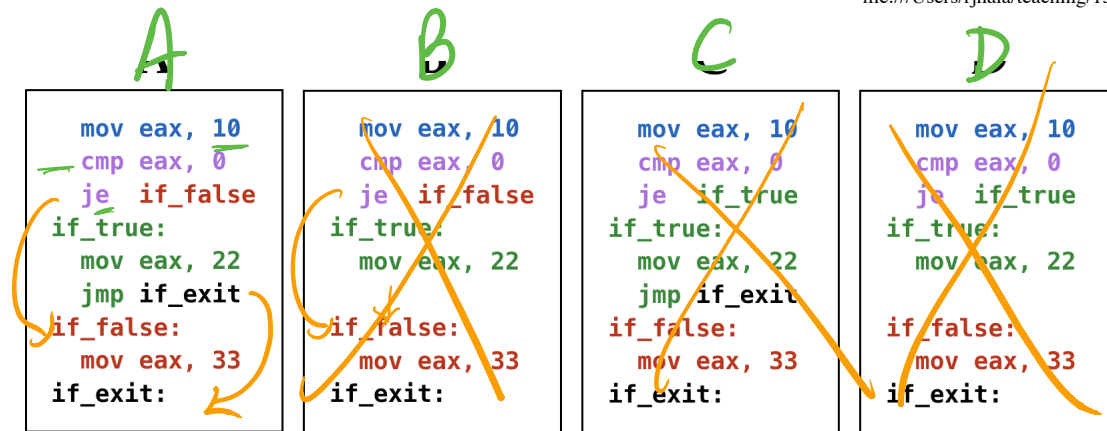
Use the result of the **flag** set by the most recent `cmp` \* To *continue execution* from the given LABEL

## QUIZ

Which of the following is a valid x86 encoding of

```
if 10:
    22
else
    33
```

**A****B****C****D**



QUIZ: Compiling if-else

## Strategy

To compile an expression of the form

```

if eCond:
  eThen
else:
  eElse

```

```

<eCond>
cmp eax, 0
je if-false
if-true:
  <eThen>
  jmp if-exit
if-false:
  <eElse>
  if-exit

```

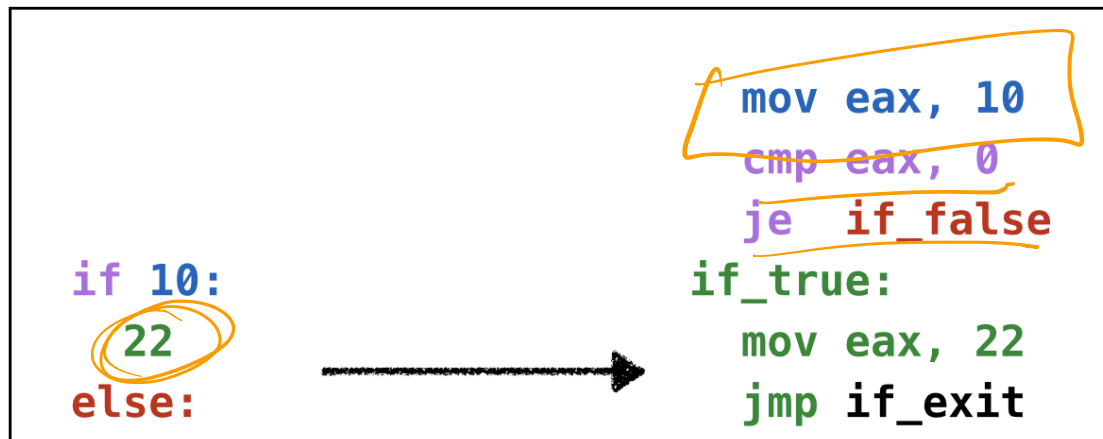
We will:

1. Compile eCond
2. Compare the result (in `eax`) against 0
3. Jump if the result is zero to a **special** "IfFalse" label
  - At which we will evaluate eElse ,
  - Ending with a special "IfExit" label.
4. (Otherwise) continue to evaluate eTrue
  - And then jump (unconditionally) to the "IfExit" label.

## Example: If-Expressions to ASM

Lets see how our strategy works by example:

### Example: if1



```
sub1(0)
```

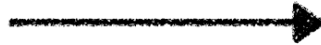
```
if_false:  
    mov eax, 0  
    sub eax, 1  
if_exit:
```

Example: if1

*Example: if2*

```
mov eax, 1  
sub eax, 1  
cmp eax, 0  
je if_false
```

```
if sub(1):  
    22  
else:  
    sub1(0)
```



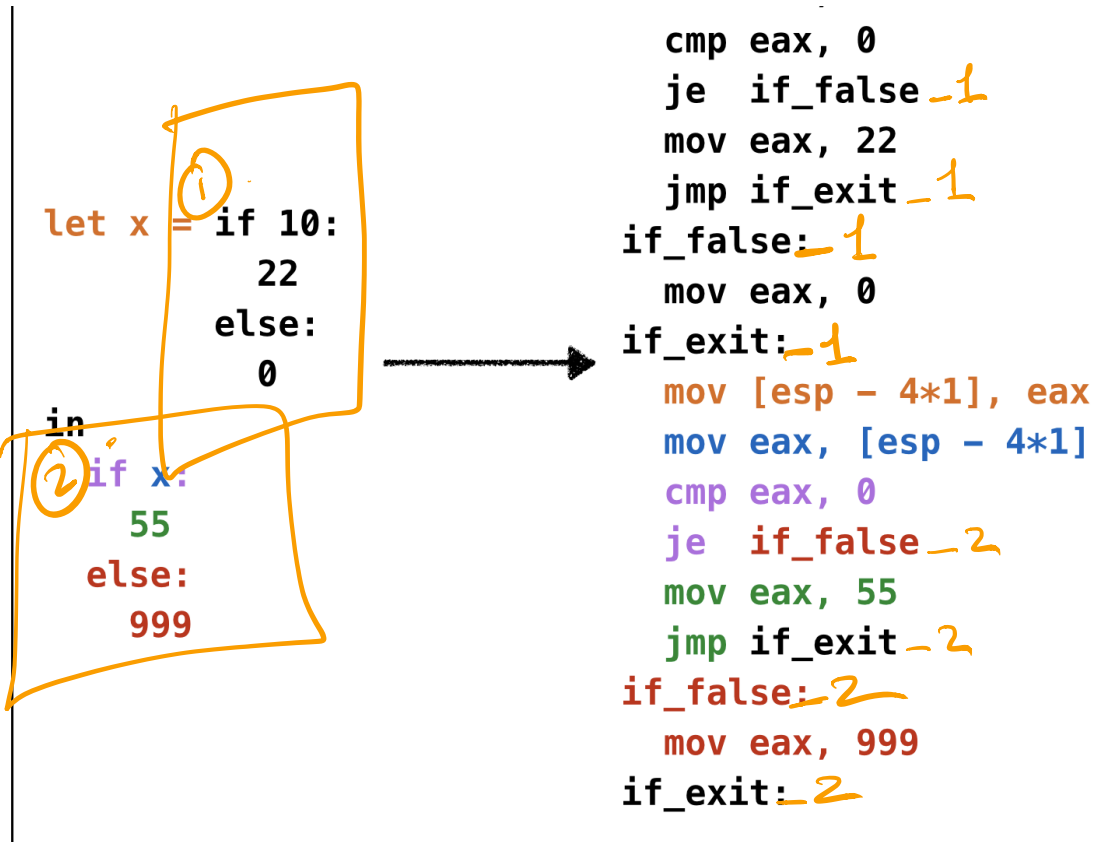
```
if_true:  
    mov eax, 22  
    jmp if_exit  
if_false:  
    mov eax, 0  
    sub eax, 1  
if_exit:
```

Example: if2

*Example: if3*

*ASM*

```
mov eax, 10
```

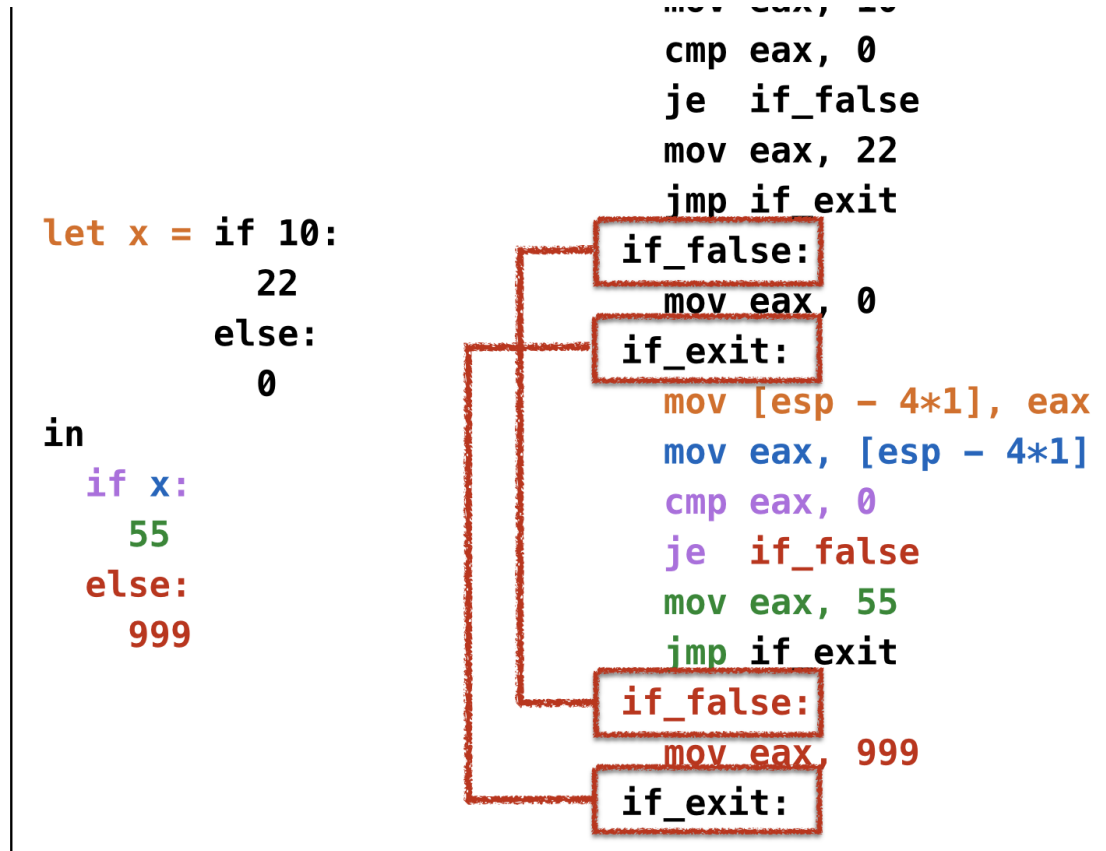


Example: if3

Oops, **cannot reuse labels** across if-expressions!

- Can't use same label in two places (invalid assembly)

```
mov eax, 10
```



Example: if3 wrong

Oops, need **distinct labels** for each branch!

- Require **distinct tags** for each if-else expression

```
mov eax, 10
```



```

let x = if 10:
  1 22
  else:
    0
in
  if x:
    2 55
  else:
    999

```

```

mov eax, 10
cmp eax, 0
je if_1_false
mov eax, 22
jmp if_1_exit
if_1_false:
mov eax, 0
if_1_exit:
mov [esp - 4*1], eax
mov eax, [esp - 4*1]
cmp eax, 0
je if_2_false
mov eax, 55
jmp if_2_exit
if_2_false:
mov eax, 999
if_2_exit:

```

Example: if3 tagged

